

A Memory-Efficient Algorithm for Large-Scale Symmetric Tridiagonal Eigenvalue Problem on Multi-GPU Systems

Hyunsu Cho and Peter A. Yoon
Trinity College, Hartford, CT, USA

Symmetric Eigenvalue Problem

$$A\mathbf{x} = \lambda\mathbf{x}$$

where A is symmetric

Many interesting applications require eigenvectors

Divide and Conquer

Yields **full spectrum** of eigenvalues and eigenvectors

Is numerically stable

Gives rise to **independent subproblems**

Often faster than $O(n^3)$ due to deflation

Divide and Conquer

Apply **orthogonal similarity transformation** to reduce A to tridiagonal form

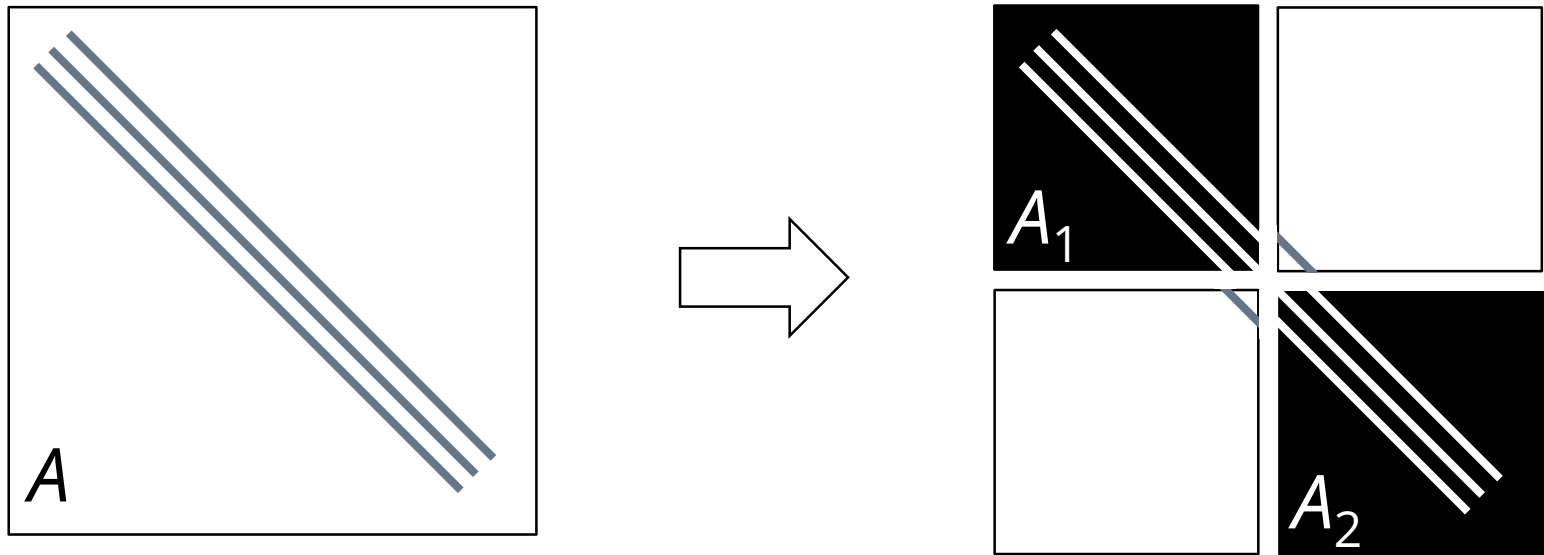
$$Q^T A Q = A'$$

where

A' is symmetric tridiagonal
and Q is orthogonal

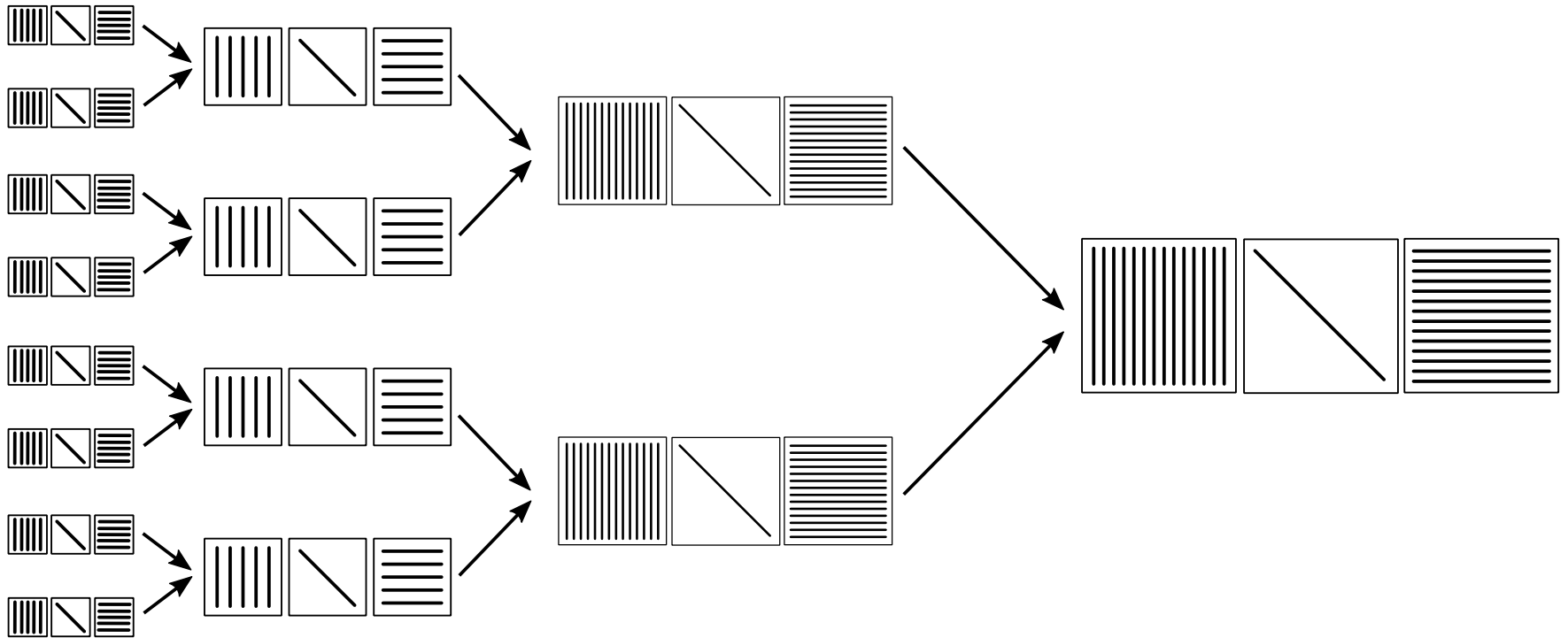
Existing work on single-node, multi-GPU:
MAGMA (UTK)

Divide and Conquer



- Solve subproblems
- Merge solutions
- Repair

Divide and Conquer



Merging solutions

Suppose

$$A = \left[\begin{array}{c|c} A_1 & \\ \hline & A_2 \end{array} \right] + \left[\begin{array}{c|c} & b_m \\ \hline b_m & b_m \end{array} \right]$$

where

$$A_1 = Q_1 D_1 Q_1^T \quad (\text{subproblem \#1})$$

$$A_2 = Q_2 D_2 Q_2^T \quad (\text{subproblem \#2})$$

Merging solutions

Then

$$A = QDQ^T + \left[\begin{array}{c|c} & \\ \hline b_m & b_m \\ \hline b_m & b_m \end{array} \right]$$

where

$$Q = \left[\begin{array}{c|c} Q_1 & \\ \hline & Q_2 \end{array} \right] \quad \text{and} \quad D = \left[\begin{array}{c|c} D_1 & \\ \hline & D_2 \end{array} \right]$$

Merging solutions

Rank-one modifier

$$H = b_m \begin{bmatrix} \mathbf{e}_m \\ \mathbf{e}_1 \end{bmatrix} \begin{bmatrix} \mathbf{e}_m \\ \mathbf{e}_1 \end{bmatrix}^T$$

Then

$$A = QDQ^T + \begin{bmatrix} & | & \\ & b_m & b_m \\ \hline & b_m & b_m \\ & | & \end{bmatrix}$$

where

$$Q = \begin{bmatrix} Q_1 & | & \\ \hline & Q_2 & \end{bmatrix} \quad \text{and} \quad D = \begin{bmatrix} D_1 & | & \\ \hline & D_2 & \end{bmatrix}$$

Rank-one update

$$H = b_m \begin{bmatrix} \mathbf{e}_m \\ \mathbf{e}_1 \end{bmatrix} \begin{bmatrix} \mathbf{e}_m \\ \mathbf{e}_1 \end{bmatrix}^T$$

$$\begin{aligned} A &= QDQ^T + H \\ &= Q(D + b_m \mathbf{z}\mathbf{z}^T)Q^T \end{aligned}$$

where

$$\mathbf{z} = Q^T \begin{bmatrix} \mathbf{e}_m \\ \mathbf{e}_1 \end{bmatrix} = \begin{bmatrix} \text{last column of } Q_1^T \\ \text{first column of } Q_2^T \end{bmatrix}$$

Rank-one update

$$H = b_m \begin{bmatrix} \mathbf{e}_m \\ \mathbf{e}_1 \end{bmatrix} \begin{bmatrix} \mathbf{e}_m \\ \mathbf{e}_1 \end{bmatrix}^T$$

$$\begin{aligned} A &= Q D Q^T + H \\ &= Q (D + b_m \mathbf{z} \mathbf{z}^T) Q^T \end{aligned}$$

where

$$\mathbf{z} = Q^T \begin{bmatrix} \mathbf{e}_m \\ \mathbf{e}_1 \end{bmatrix} = \begin{bmatrix} \text{last column of } Q_1^T \\ \text{first column of } Q_2^T \end{bmatrix}$$

Need eigen-decomposition of inner system

Decompose $D + b_m \mathbf{z}\mathbf{z}^T$

1. Sort entries in D ; permute \mathbf{z} likewise
2. Filter some entries in D and \mathbf{z} via deflation (next slide)

Decompose $D + b_m \mathbf{z}\mathbf{z}^T$

1. Sort entries in D ; permute \mathbf{z} likewise
2. Filter some entries in D and \mathbf{z} via deflation (next slide)
3. Compute all roots of the **secular equation** [1]

$$1 + b_m \sum_{i=1}^n \frac{d_i^2}{z_i - \lambda} = 0,$$

giving the m eigenvalues.

4. Compute corresponding eigenvectors stably [2]

[1] Li 1994

[2] Gu & Eisenstat 1994

Decompose $D + b_m \mathbf{z}\mathbf{z}^T$

1. Sort entries in D ; permute \mathbf{z} likewise
2. Filter some entries in D and \mathbf{z} via deflation (next slide)
3. Compute all roots of the **secular equation** [1]

$$1 + b_m \sum_{i=1}^n \frac{d_i^2}{z_i - \lambda} = 0,$$

giving the m eigenvalues.

4. Compute corresponding eigenvectors stably [2]
5. Multiply each eigenvector by Q

Recall: $A = Q(D + b_m \mathbf{z}\mathbf{z}^T)Q^T$

[1] Li 1994

[2] Gu & Eisenstat 1994

Deflation

Recall:

$$D = \left[\begin{array}{c|c} D_1 & \\ \hline & D_2 \end{array} \right]$$

Entries of D are eigenvalues of two subproblems

If two entries are nearly identical, we throw one away

Fewer columns when multiplying eigenvectors by Q

Same thing for small entries in z

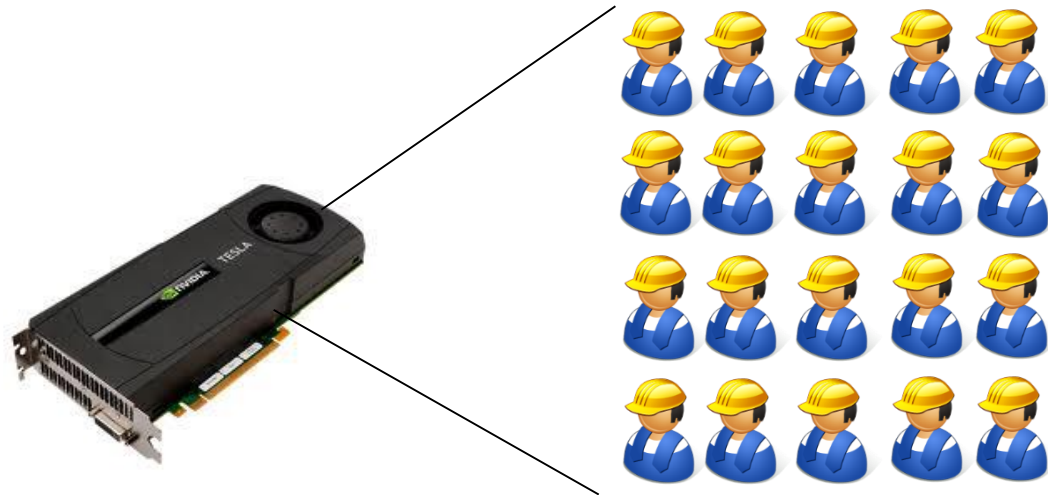
Reduce work complexity to $O(n^{2.3})$

GPU computing

General-purpose computation on GPUs

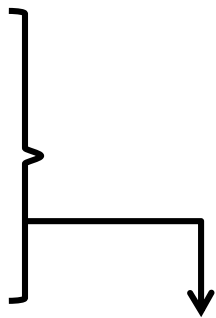
Bulk parallelism w/ many small threads

Cost effective; widely available



Mapping work to GPU

1. Sort entries in D ; permute z likewise
2. Filter some entries in D and z via deflation
3. Compute all roots of the secular equation, giving the m eigenvalues.
4. Compute corresponding eigenvectors stably
5. **Multiply each eigenvector by Q**
→ Done in bulk via DGEMM



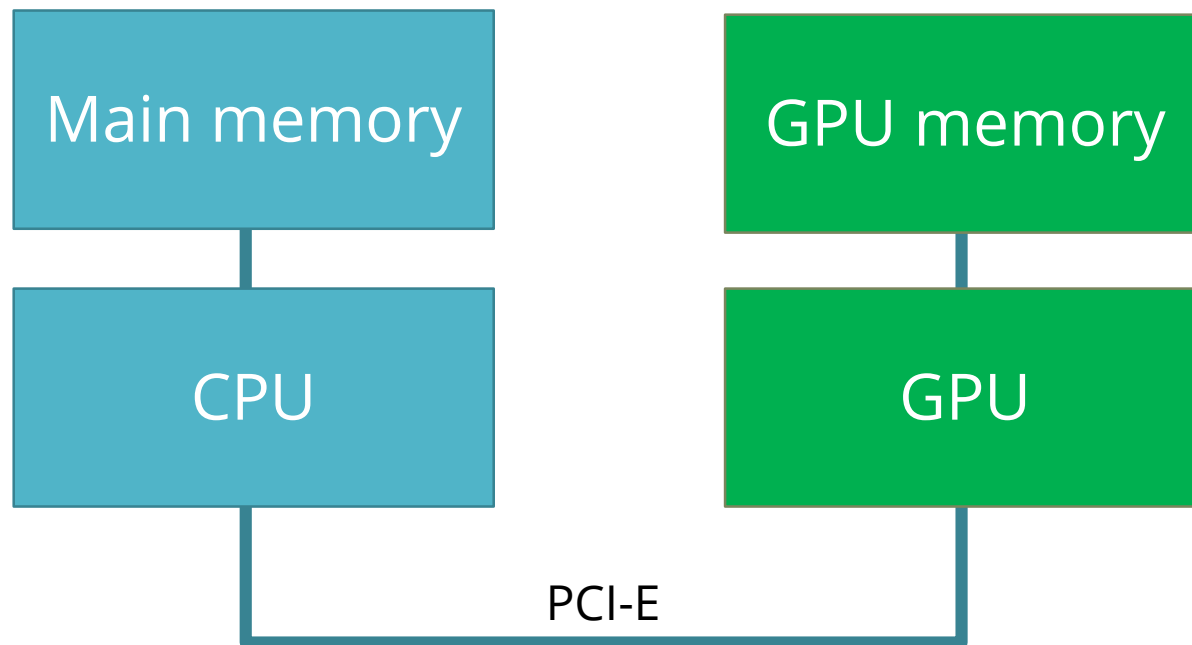
Parallel but
not as work-intense

GPU memory

High-bandwidth dedicated memory

Separate from main memory

Limited in size



Memory requirement

Eigenvectors are dense

→ $O(n^2)$ storage

Intermediate workspace: eigenvectors of inner system

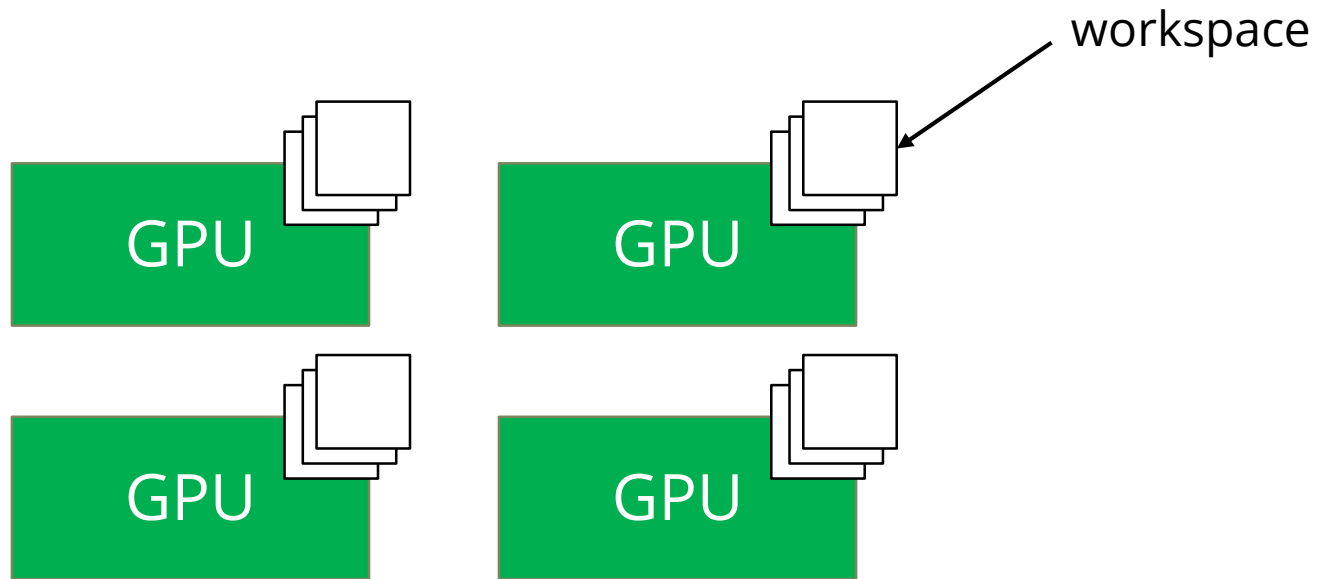
| Matrix dimension | Memory required |
|------------------|-----------------|
| 8192 | 1.5 GB |
| 16384 | 5.8 GB |
| 32768 | 23.4 GB |
| 36000 | 28.2 GB |
| 50000 | 54.4 GB |

Our contribution

Overcome limitation in GPU memory
while retaining adequate performance

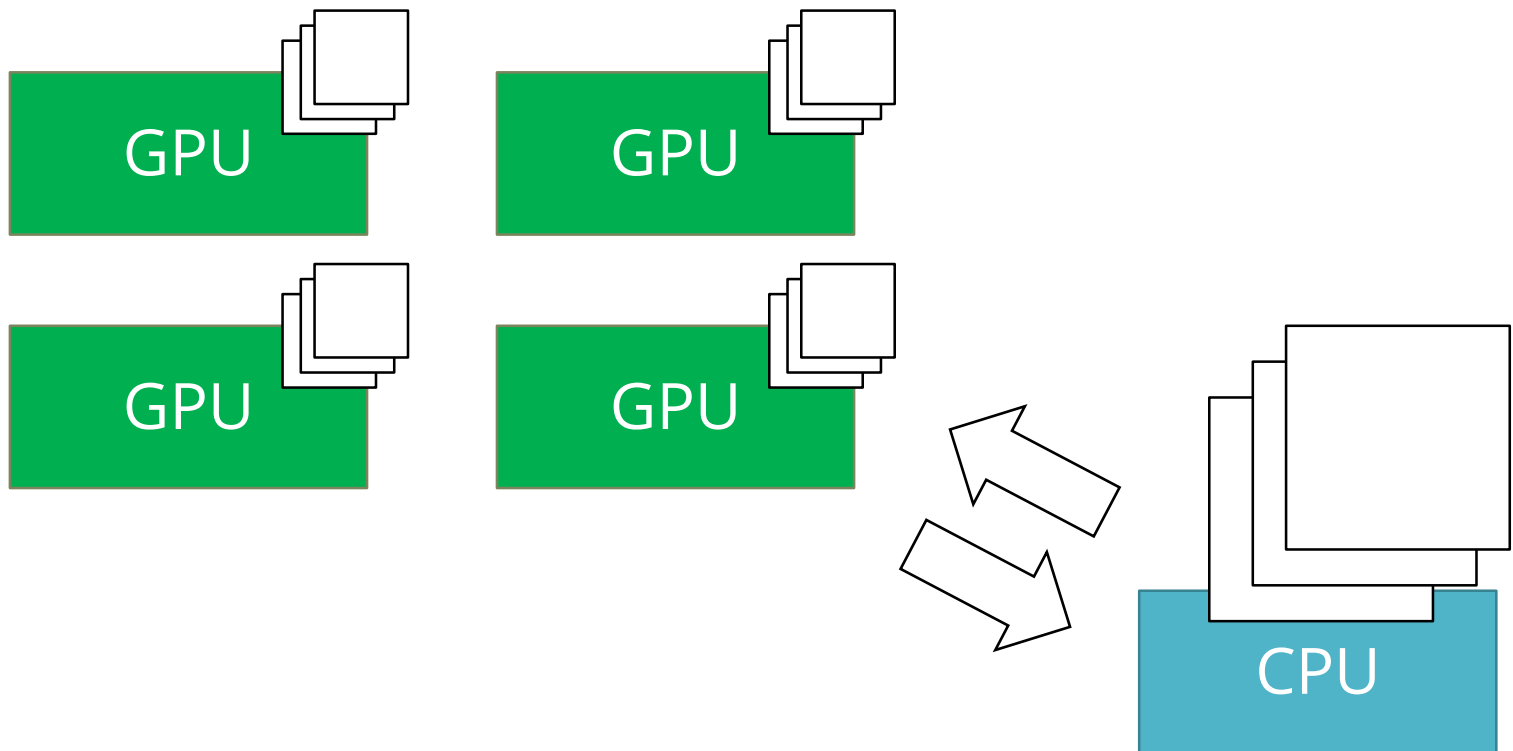
Strategies

1. Use multiple GPUs



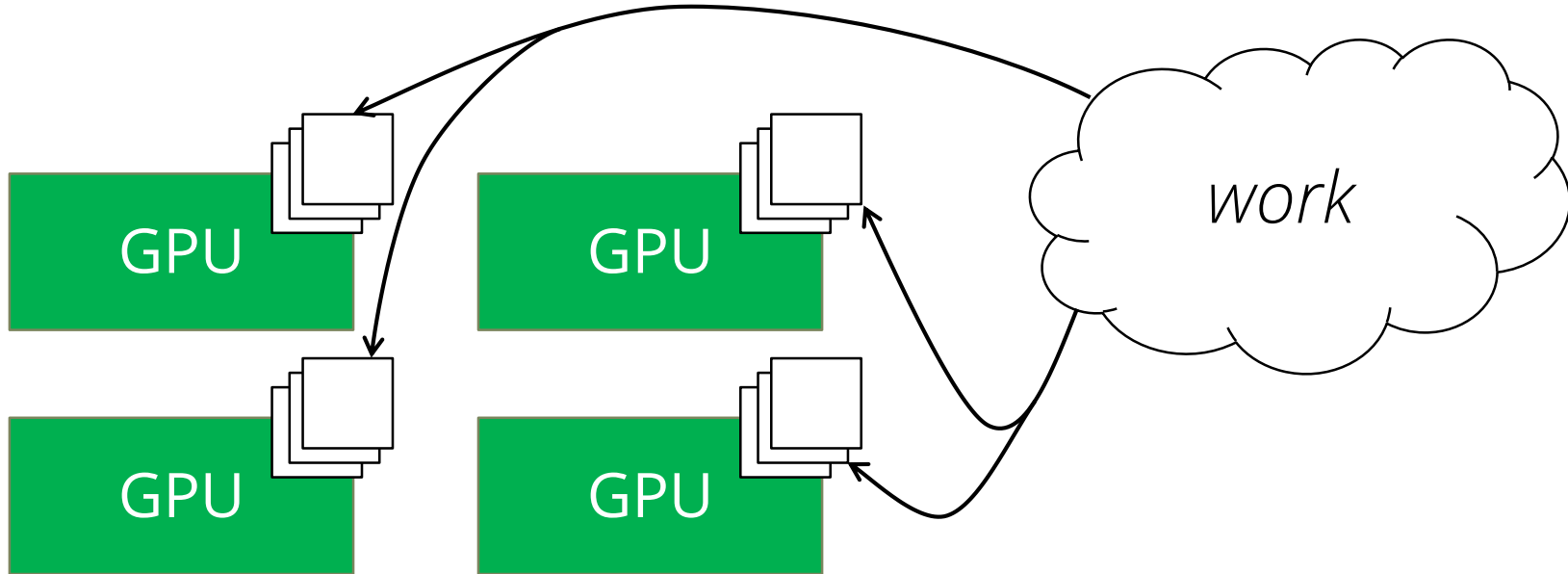
Strategies

2. Keep most of workspace in main memory
(**out-of-core** approach)



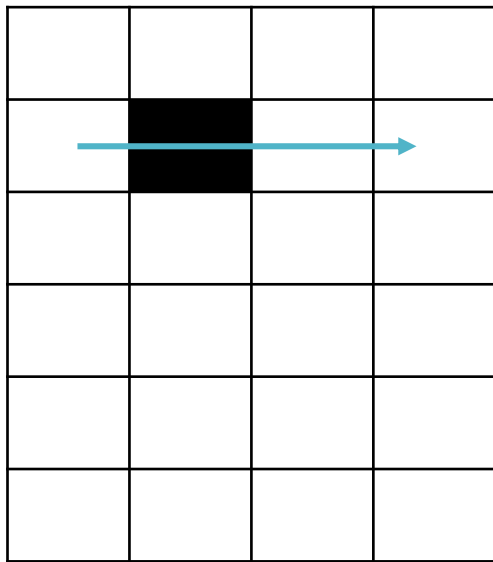
Strategies

3. **Shape work** to fit GPU workspaces



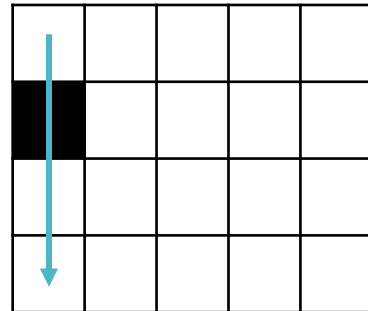
Block matrix multiplication

Use a fine partition to fit submatrices into GPU memory

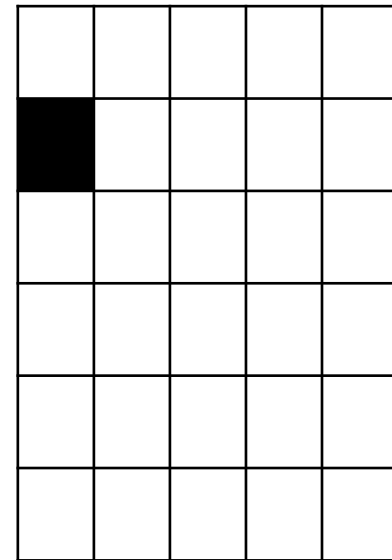


Q

X



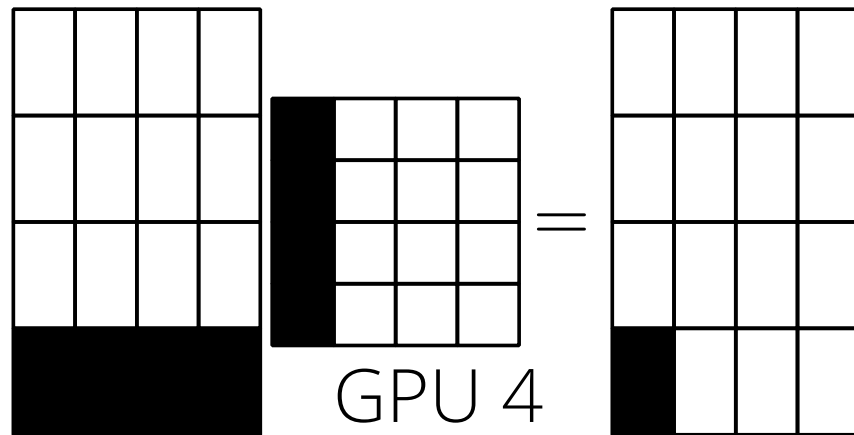
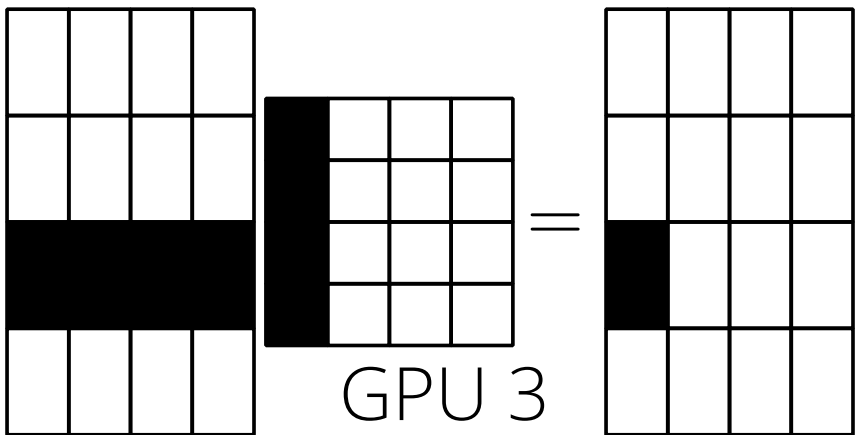
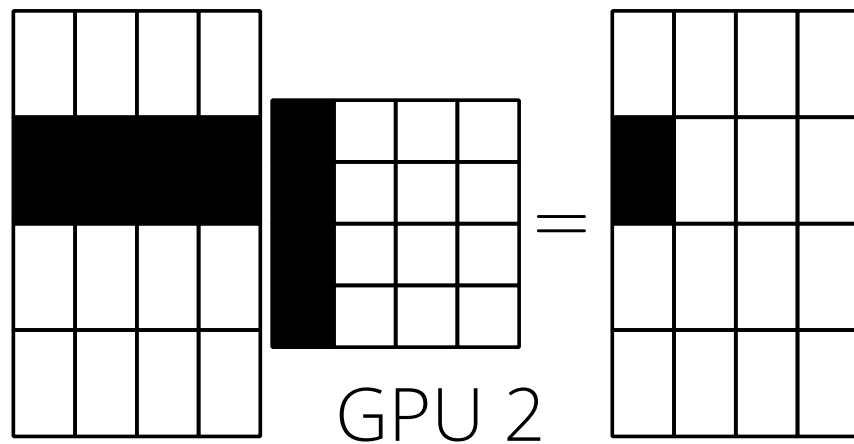
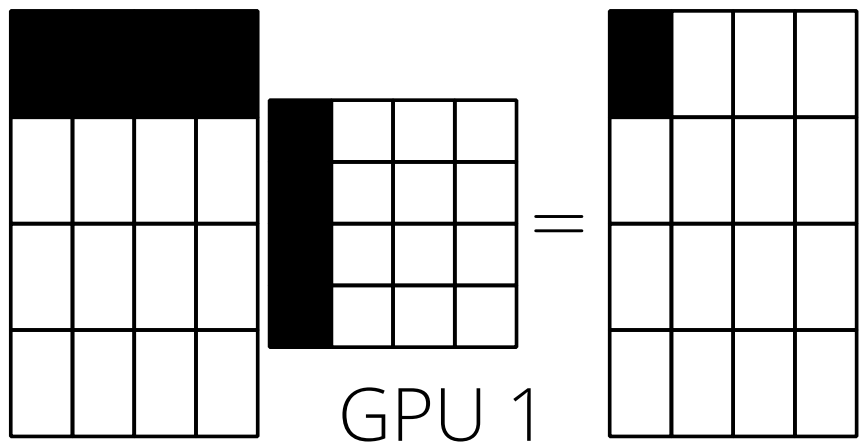
=



Eigenvectors of

$$D + b_m \mathbf{z}\mathbf{z}^T$$

Eigenvectors of A

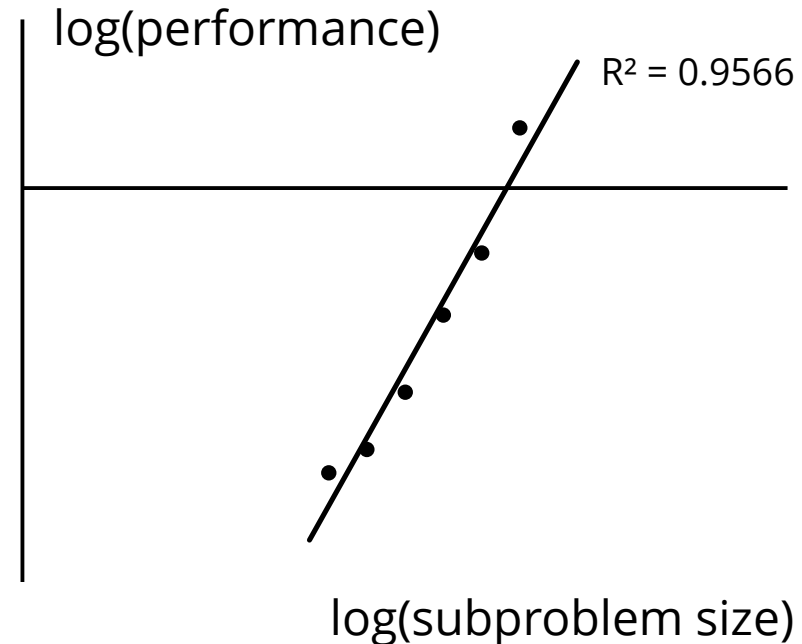
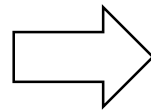
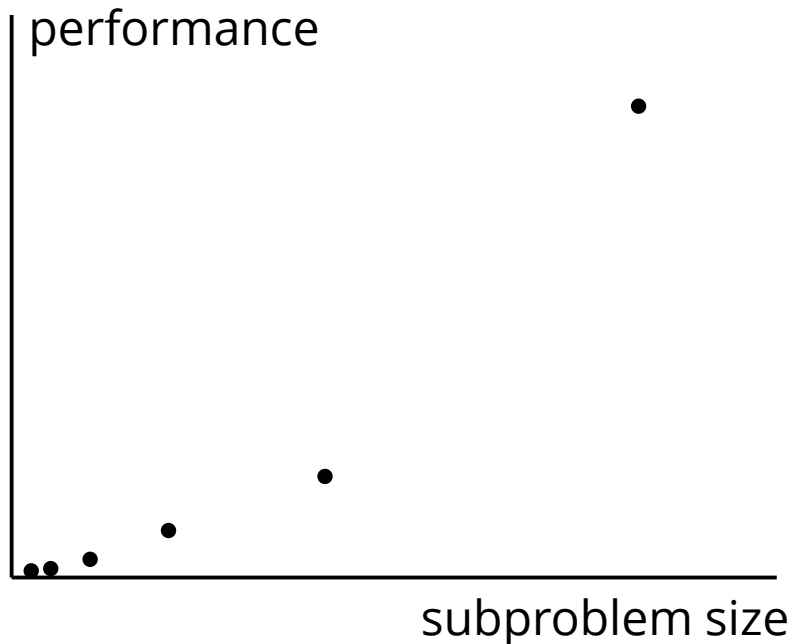


Hybrid computation

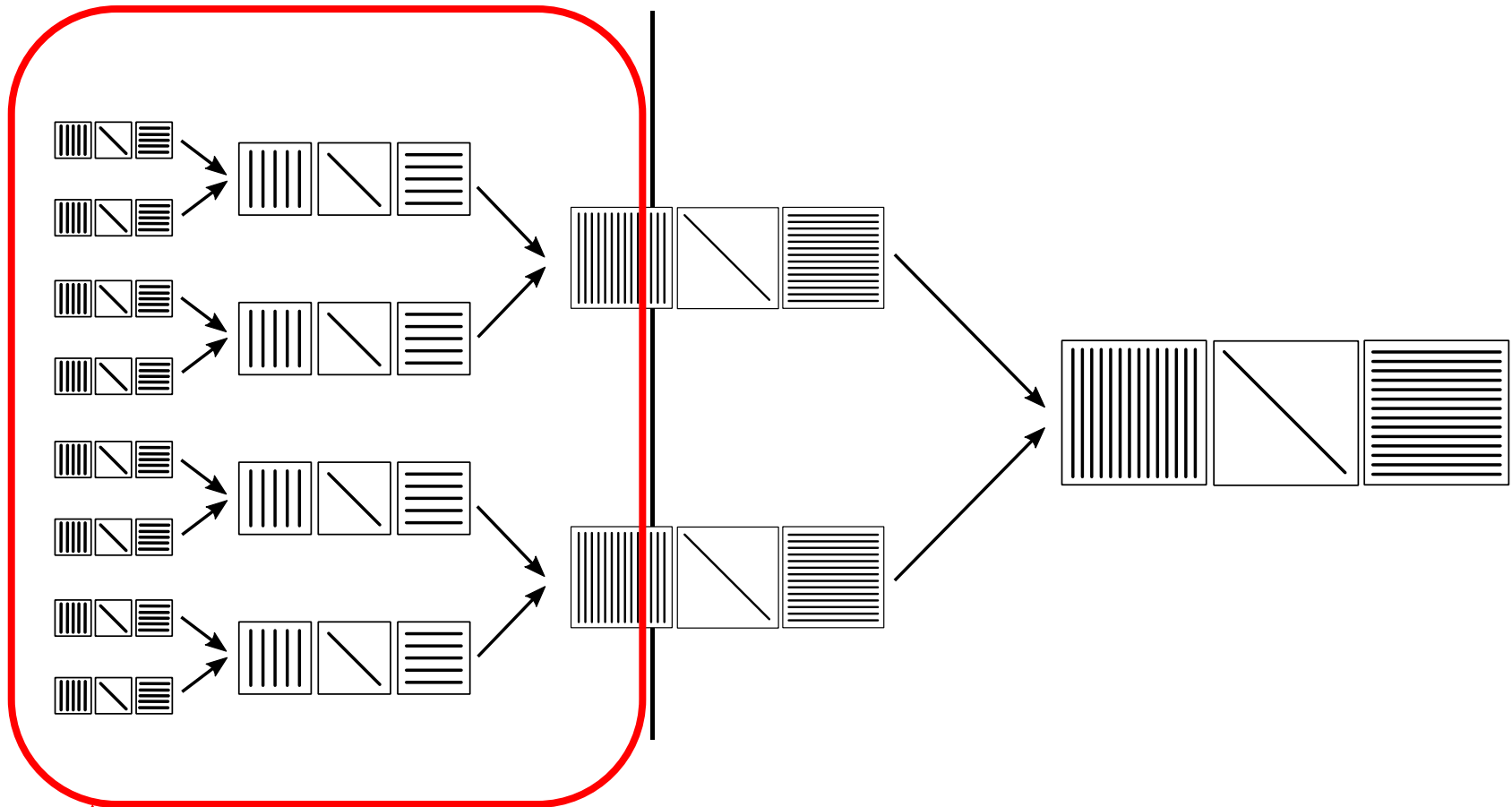
Allocate subproblems to both GPUs and CPUs

Model performance as a power function

Profiler fits parameters using least-squares

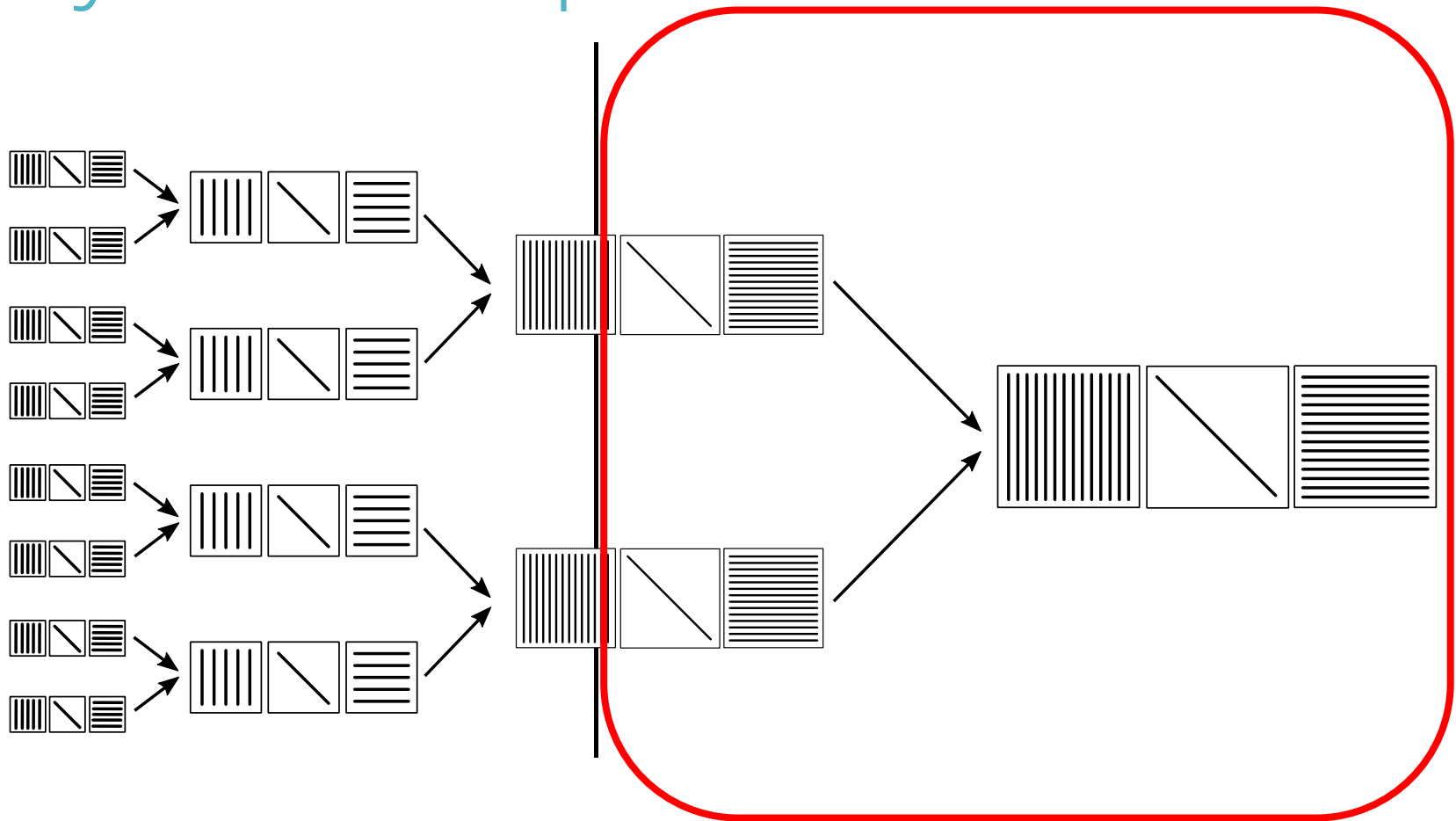


Hybrid computation



↑ Solve many subproblems in parallel

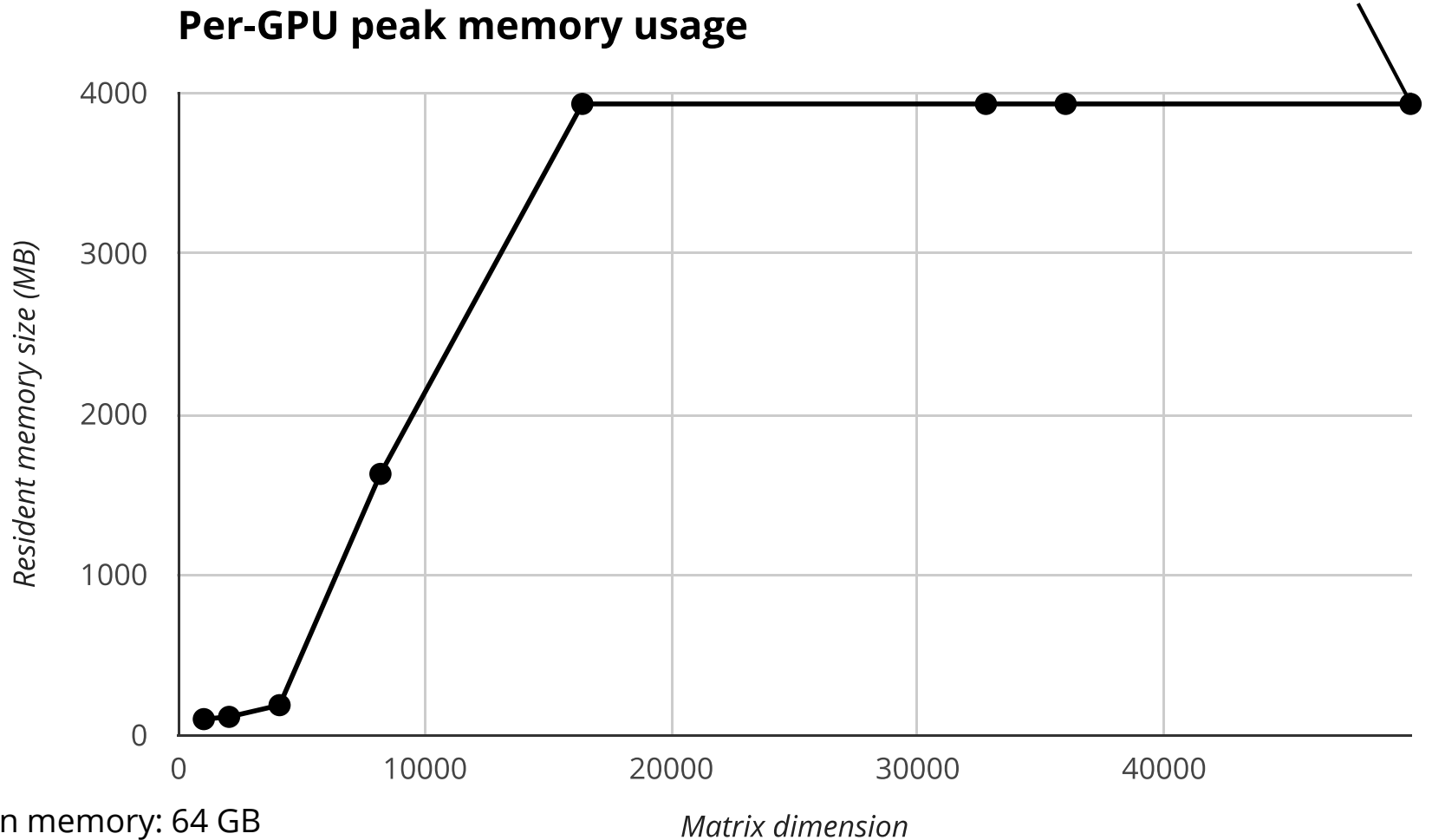
Hybrid computation



Solve each subproblem by parts \longrightarrow

Results

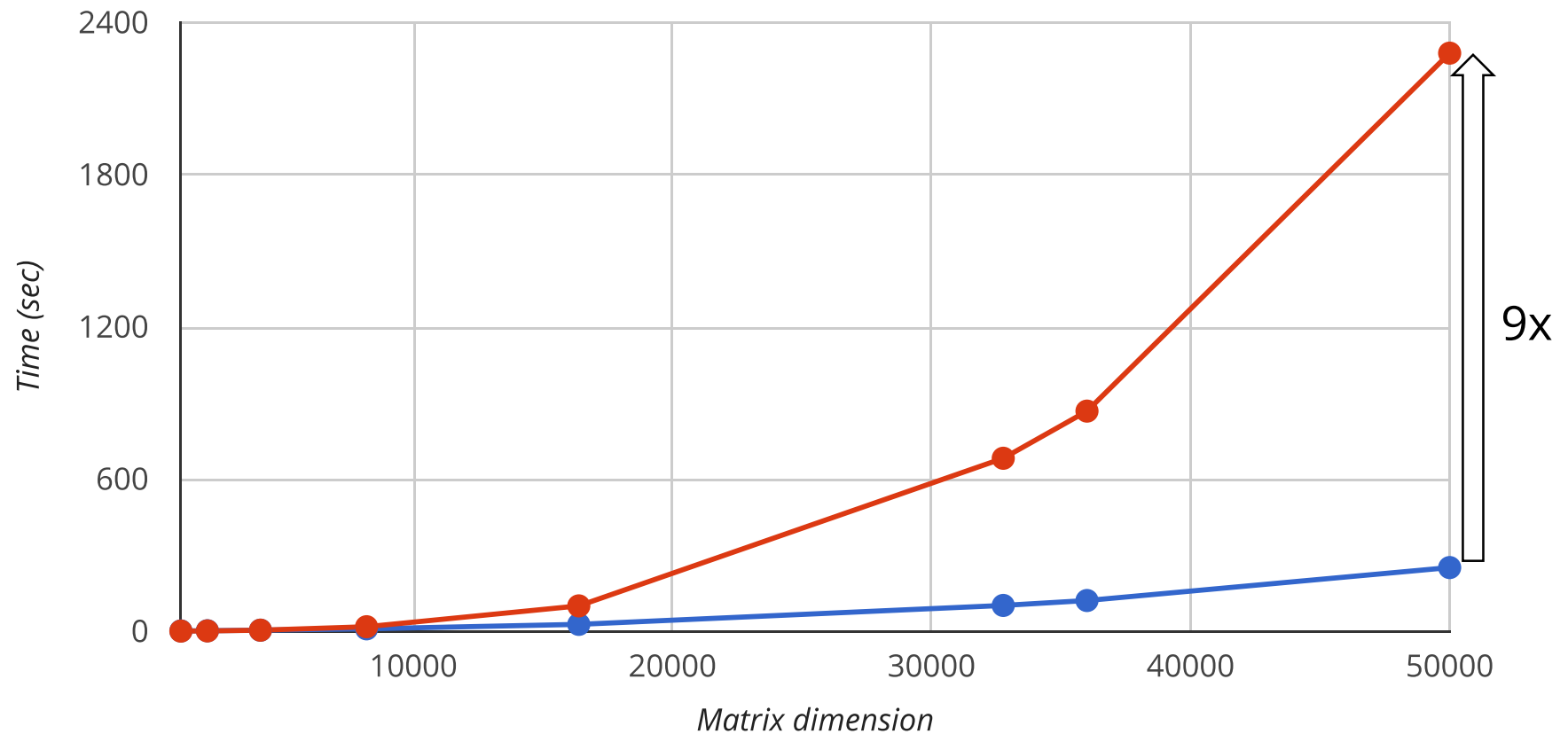
Scales to 50k * 50k matrix
With 4 GB of GPU memory



Main memory: 64 GB
GPU memory: 5 GB per GPU

Results

Performance: vs. multicore CPU



CPU: dual Intel® Xeon® E5-2620
GPU: 4 Nvidia Tesla® K20c

■ GPUs + CPUs ■ CPUs only

Conclusion

Out-of-core approach overcomes memory limitation on the GPU

Hybrid computation with profiling delivers reasonable performance

Acknowledgment

Trinity College, Student Research Program

Nvidia Corporation, CUDA Teaching Center Program

Any questions?